## *Polynomials in Action*<sup>1</sup>

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We review several examples of how polynomials interact with monoid and category actions. Written at the 2024 Poly at Work Workshop at the Topos Institute.

## 1 Dynamical systems with actions

To discuss continuous-time dynamics in **Poly**, we need some notion of an  $\mathbb{R}$ -action. More generally, we would like to consider how dynamical systems in **Poly** may be equipped with the action of a monoid (M, e, \*) (in **Set**), another polynomial, or perhaps any category. Here are several ways we may model such a concept with polynomials.

- 1. We could consider cofunctors  $\mathbb{S} \rightarrow y^M \otimes \mathbb{C}_p$ , where  $y^M$  is the carrier of the polynomial comonad corresponding to M viewed as a 1-object category, while  $\mathbb{C}_p$  is the cofree comonad on a polynomial functor p. We could replace  $y^M$  with an arbitrary category. We could also take an appropriate subcategory of  $y^M \otimes \mathbb{C}_p$ , associating morphisms in  $\mathbb{C}_p$  (tuples of directions in p) with appropriate time values in M.
- We could consider lenses φ: Sy<sup>S</sup> → [My, p] such that, for every section γ: p → y, composing φ with [My, γ] yields a cofunctor Sy<sup>S</sup> → [My, y] ≅ y<sup>M</sup>.<sup>2</sup> We could also consider lenses Sy<sup>S</sup> ⊗ My → p that become cofunctors when composed with any section p → y.
- 3. We could consider *p*-coalgebras for polynomial endofunctors *p* on  $[y^M, \mathbf{Set}]$ . See the next section for more exposition on this.
- 4. We could consider the category of (y<sup>M</sup>, y<sup>M</sup>)-bicomodules, the category of (C<sub>p</sub>, C<sub>p</sub>)-bicomodules, or perhaps most generally the category of (C, C)-bicomodules; and the morphisms there whose domain is a comonoid. It turns out that a (C, C)-bicomodules comonad carried by a polynomial D corresponds to a cofunctor D → C (in particular, D itself must be a polynomial comonad).
- 5. Let *S* be a smooth manifold and  $f: \mathbb{R} \to S$  be a smooth map. In the thin double category whose objects are polynomials, vertical arrows are lenses, and horizontal arrows are charts,<sup>3</sup> we can

<sup>1</sup> Inspired by conversations with Matteo Capucci, Harrison Grodin, Sophie Libkind, Toby Smithe, and David Spivak.

<sup>2</sup> This idea comes from Smithe's Open Dynamical Systems as Coalgebras for Polynomial Functors, with Application to Predictive Processing (2022), Definition 2.1.

<sup>3</sup> Introduced by David Jaz Myers in Double Categories of Open Dynamical Systems (2020). consider squares



The right lens is our open dynamical system with state space S and interface p. The left lens is vertical and picks out unit tangent vectors on directions. The bottom chart can be considered a trajectory of position and direction pairs in p.<sup>4</sup>

## 2 Polynomial functors over sets with monoid actions

Let (M, e, \*) be a monoid (in **Set**). Its associated 1-object category, viewed as a comonoid object in (**Poly**, y,  $\triangleleft$ ), is carried by the polynomial  $y^M$ , which we will use to denote this category. Then an *M*-set is a functor  $X: y^M \rightarrow$ **Set**, or a set X with an M-action  $\cdot: M \times X \rightarrow X$  respecting e and \*. A morphism of M-sets  $X \rightarrow Y$  is a natural transformation from X to Y as functors  $y^M \rightarrow$ **Set**, or an *M*-equivariant map  $f: X \rightarrow Y$ , satisfying  $f(m \cdot x) = m \cdot f(x)$  for all  $m \in M$  and  $x \in X$ . In other words, the category of M-sets and M-equivariant maps can be identified with the functor category  $[y^M,$ **Set**].

Following Gambino-Kock, we characterize polynomial endofunctors on  $[y^M, \mathbf{Set}]$ .<sup>5</sup> As a presheaf category,  $[y^M, \mathbf{Set}]$  is a topos. In particular, it is complete and locally cartesian closed: given an *M*equivariant map  $p: E \to B$ , the functor between slice categories

$$\Delta_p: [y^M, \mathbf{Set}]/B \to [y^M, \mathbf{Set}]/E$$

induced by pullback along *p* has a right adjoint

$$\Pi_p: [y^M, \mathbf{Set}]/E \to [y^M, \mathbf{Set}]/B.$$

Composing this on one side with the product functor  $\Delta_! : [y^M, \mathbf{Set}] \rightarrow [y^M, \mathbf{Set}]/E$  sending  $X \mapsto (E \times X \xrightarrow{\pi} X)$  and on the other with the forgetful functor  $\Sigma_! : [y^M, \mathbf{Set}]/B \rightarrow [y^M, \mathbf{Set}]$  sending  $(Y \rightarrow B) \mapsto Y$  yields the polynomial functor corresponding to  $p : E \rightarrow B$ :

$$\Sigma_!\Pi_p\Delta_!\colon [y^M, \mathbf{Set}] \to [y^M, \mathbf{Set}],$$

which we will also denote by p. Then the category of polynomial endofunctors on  $[y^M, \mathbf{Set}]$ , denoted  $\mathbf{Poly}_{[y^M, \mathbf{Set}]}$ , has these polynomial objects as functors and all natural transformations<sup>6</sup> between them as morphisms. Equivalently (and we will freely switch between the two characterizations), it is the category where <sup>4</sup> I learned this from Matteo Capucci.

<sup>5</sup> See *Polynomial Functors and Polynomial Monads* (2009) by Gambino and Kock; we follow their notation. All this can be generalized to an arbitrary locally cartesian closed category (with a terminal object and pullbacks).

<sup>6</sup> Gambino-Kock consider only the *cartesian* natural transformations, but we would like to consider all of them.

- an object is an *M*-equivariant map  $p: E \rightarrow B$ ;
- a morphism  $\varphi$  from  $p: E \to B$  to  $p': E' \to B'$  consists of
  - an *M*-equivariant map  $\varphi_1 \colon B \to B'$ ;
  - an *M*-equivariant map  $\varphi^{\sharp} \colon E' \times_{B'} B \to E$ , whose domain is the pullback of p' along  $\varphi_1$ .

There is a functor  $\operatorname{Poly} \to \operatorname{Poly}_{[y^M, \operatorname{Set}]}$  induced by the functor  $\operatorname{Set} \to [y^M, \operatorname{Set}]$  that sends every set to its constant presheaf, i.e. the same set with a trivial action. There is also a functor  $\operatorname{Poly}_{[y^M, \operatorname{Set}]} \to$   $\operatorname{Poly}$  that forgets *M*-actions. Each polynomial in  $\operatorname{Poly}_{[y^M, \operatorname{Set}]}$  therefore has an underlying polynomial  $\sum_{i \in I} \prod_{a \in A_i} y$ , along with an *M*action on *I* and a compatible *M*-action on  $\sum_{i \in I} A_i$ .<sup>7</sup>

We can lift many<sup>8</sup> of the usual structures on **Poly** to this setting. In particular, given  $p \in \mathbf{Poly}_{[y^M, \mathbf{Set}]}$ , we can consider *p*-coalgebras  $S \to p(S)$  for  $S \in [y^M, \mathbf{Set}]$ , or perhaps<sup>9</sup> equivalently morphisms from  $S \times S \to S$  to *p* in  $\mathbf{Poly}_{[y^M, \mathbf{Set}]}$ . In **Poly**, these are our open dynamical systems. Taking  $M := \mathbb{R}$  may then be a way to discuss open dynamical systems with some notion of time  $t \in \mathbb{R}$ . <sup>7</sup> Given  $i \in I$  and  $a \in A_i$ , the *M*-action must send (i, a) to  $(m \cdot i, b)$  for some  $b \in A_{m \cdot i}$ .

<sup>8</sup> It would be interesting to verify precisely which structures on **Poly** generalize and which do not.
<sup>9</sup> This needs verification (or perhaps it simply isn't immediately obvious to me).